

Directed Algebraic Topology, and applications to static analysis of programs

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Modelisation and Analysis of Systems in Interaction lab

CEA/Saclay and Ecole Polytechnique

Outline of the talk

- ▶ A very short introduction to static analysis, in view of verifying concurrent programs
- ▶ Reducing the state-space model: components categories
 - ▶ “Good properties” of component categories: lifting property and van Kampen
- ▶ A closer look at components: future and past components
 - ▶ lifting and van Kampen
 - ▶ reflectivity in $\overrightarrow{\pi_1}(\overrightarrow{X})$
- ▶ Some mathematical issues

Context

Static analysis of programs

- ▶ Find **outer-approximation** of sets of reachable values of variables at some program points
- ▶ To ensure **absence of runtime errors** typically

Example

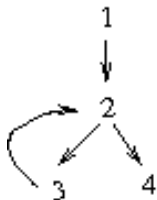
```
float x;
```

```
x=[0,1];           [1]
```

```
while (x<=1) {    [2]
```

```
    x = x-0.5*x;  [3]
```

```
}                [4]
```



$$x_1 = [0, 1]$$

$$x_2 =]-\infty, 1] \cap (x_1 \cup x_3)$$

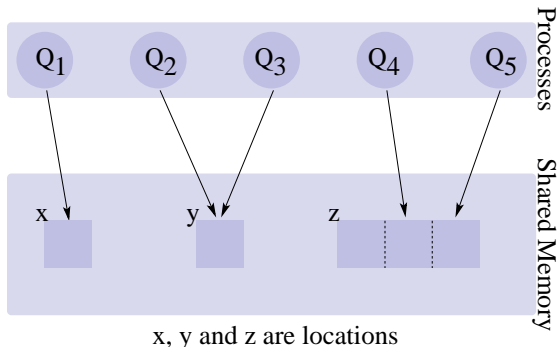
$$x_3 = x_2 - 0.5x_2$$

$$x_4 =]1, \infty[\cap x_2$$

(final **smallest** invariant: $x_2 \in [0, 1]$, $x_4 = \emptyset$)

Concurrent programs

shared memory style



Not sequential programs, bad states, chaotic behavior

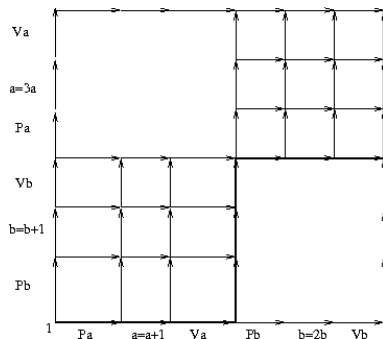
\implies Need for synchronizations \implies Need for locks: P_y , V_y

\implies Interleaving semantics given by a “shuffle” of transition systems
(or fibred product)

Interleaving semantics of...

$$(a = 1, b = 2) Pa.a + 1.Va.Pb.2b.Vb \mid Pb.b + 1.Vb.Pa.3a.Va$$

T_1	T_2
Pa	—
$a = a + 1$	—
Va	
—	Pb
—	$b = b + 1$
—	Vb
Pb	—
$b = 2 * b$	—
Vb	—
—	Pa
—	$a = 3 * a$
—	Va



Equations (invariants)...

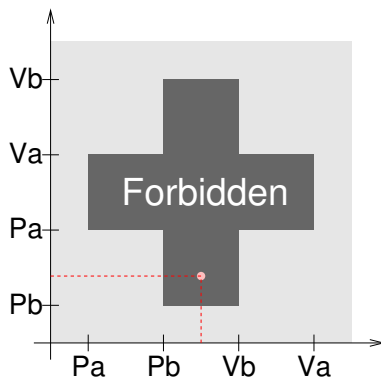
$$\left\{ \begin{array}{l} X_1 = (a_0, b_0) \\ X_2 = \dots \\ \quad = \dots \text{you don't want to know...} \\ X_{41} = \dots \end{array} \right.$$

(41 vertices, 60 edges!)

Geometry

“progress graphs” E.W.Dijkstra’68 (later V.Pratt, R. van Glabbeek’91)

$T1=Pa.Pb.Vb.Va$ in parallel with $T2=Pb.Pa.Va.Vb$



“Continuous model”: $x_i =$ local time; dark grey region=**forbidden!**

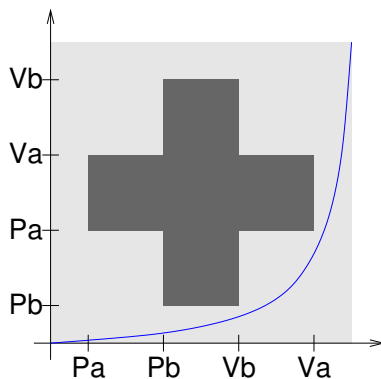
see *Algebraic Topology and Concurrency* TCS 2006, L. Fajstrup, E. Goubault, M.



Rausen

Execution paths

are continuous

$T1=Pa.Pb.Vb.Va$ in parallel with $T2=Pb.Pa.Va.Vb$



Traces are continuous paths increasing in each coordinate: [dipaths](#)  

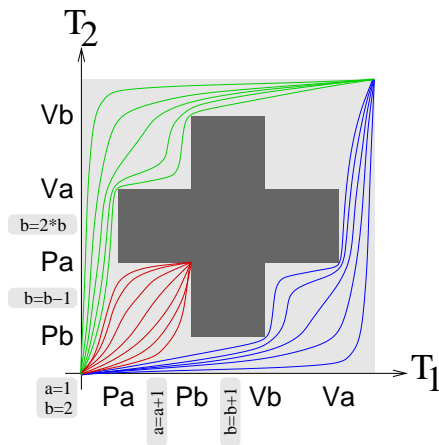
Classes of equivalent dipaths

up to dihomotopy

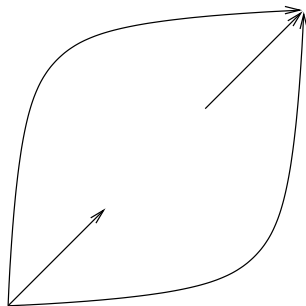
T_1 gets a and b before $T_2 \Rightarrow a=2$ and $b=4$

T_2 gets b and a before $T_1 \Rightarrow a=2$ and $b=3$

Each of T_1 and T_2 gets a resource
 \Rightarrow Deadlock with $a=2$ and $b=1$



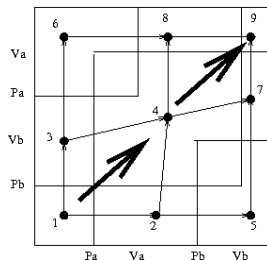
Ideally, we want to retract to...
(not quite true though)



We will get back to this later.

Use for our first example

$Pa.a + 1.Va.Pb.2b.Vb \mid Pb.b + 1.Vb.Pa.3a.Va$

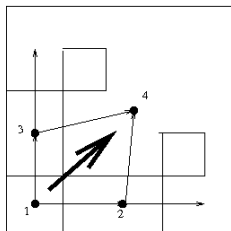


$$\left\{ \begin{array}{l} X_1 = (a_0, b_0) \\ X_2 = (a_1 + 1, b_1) \\ X_3 = (a_1, b_1 + 1) \\ X_4 = (a_1 + 1, b_1 + 1) \\ X_5 = (a_2, 2 * b_2) \\ X_6 = (3 * a_3, b_3) \\ X_7 = (a_4, 2 * b_4) \cup (a_5, b_5 + 1) \\ X_8 = (3 * a_4, b_4) \cup (a_6 + 1, b_6) \\ X_9 = (3 * a_7, b_7) \cup (a_8, 2 * b_8) \end{array} \right.$$

component category - finite number of objects, morphisms and relations!

In fact...aim of this talk is to go even further...

- ▶ They are forward equations so we need only the following retract (“future components”):



$$\left\{ \begin{array}{l} X_1 = (a_0, b_0) \\ X_2 = (a_1 + 1, b_1) \\ X_3 = (a_1, b_1 + 1) \\ X_4 = (3 * (a_1 + 1), 2 * (b_1 + 1)) \\ \quad \cup (3 * a_2, 2 * b_2 + 1) \\ \quad \cup (3 * a_3 + 1, 2 * b_3) \end{array} \right.$$

- ▶ (a similar method exists for backwards equations)
- ▶ In general, there are loops... out of the scope of this talk!

Models

- ▶ **Cubical sets** (pre-existing the field of course!)
- ▶ **Po-spaces** (i.e. topological space with closed partial order), introduced first in other fields (domain theory P. Johnstone etc., functional analysis L. Nachbin etc.) **local po-spaces** (atlas of po-spaces - L. Fajstrup, E. Goubault, M. Raussen)
- ▶ **d-spaces** (M. Grandis)
- ▶ **Flows** (P. Gaucher)
- ▶ **Streams** (S. Krishnan)
- ▶ etc.

Most of the rest would apply to all of these models (except for loops!)

Partially Ordered Spaces

framework for “progress graphs” (one only needs *MFPS'98*)

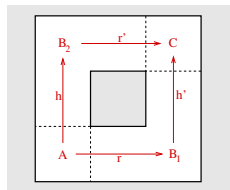
A topological space X with a (global) closed partial order \sqsubseteq

- ▶ Morphisms are increasing and continuous maps: **dimaps**
- ▶ (Finite) Traces on (X, \sqsubseteq) are dimaps from $\vec{I} = ([0, 1], \leq)$ to (X, \sqsubseteq) : **dipaths**
- ▶ **Dihomotopies** between dipaths α and β with fixed extremities x and y are dimaps $H : \vec{I} \times \vec{I} \rightarrow X$ such that for all $s \in \vec{I}$, $t \in \vec{I}$,
 - ▶ $H(t, 0) = \alpha(t)$ and $H(t, 1) = \beta(t)$
 - ▶ $H(0, s) = x$ and $H(1, s) = y$
- ▶ Two dipaths are **dihomotopic** if there exists a finite sequence of dihomotopies relating them (alternative definition, equivalent under some conditions: $H : \vec{I} \times I \rightarrow X$)

How to retract?

The fundamental category $\vec{\pi}_1(\vec{X})$ of a pospace \vec{X}

- ▶ Starting with a **variation on the Poincaré groupoid**, $\pi_1(X)$ defined as the category:
 - ▶ objects: points of X ,
 - ▶ morphisms: classes of dipaths up to dihomotopy:
a morphism from x to y is a dihomotopy class $[\alpha]$ of a dipath α going from x to y .
- ▶ We see that in most interesting (to static analysis) case, it is “**essentially**” finite



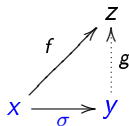
\implies Formally invert “inessential” arrows

Yoneda morphism

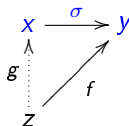
axiomatizing the preservation of the future and the past (1)

Let \mathcal{C} be a small category. A *Yoneda* morphism σ is an element of $\mathcal{C}[x, y]$ such that for all object z of \mathcal{C} ,

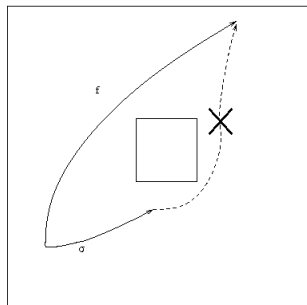
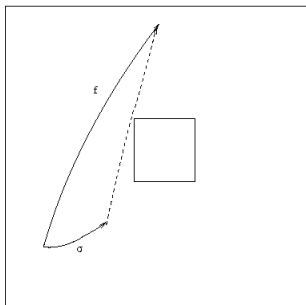
future if $\mathcal{C}[y, z] \neq \emptyset$ then for all $f \in \mathcal{C}[x, z]$, there is a unique $g \in \mathcal{C}[y, z]$ such that



past if $\mathcal{C}[z, x] \neq \emptyset$ then for all $f \in \mathcal{C}[z, y]$, there is a unique $g \in \mathcal{C}[z, x]$ such that



Example: Yoneda morphism is a “small” move

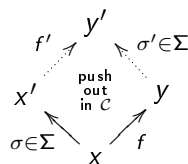
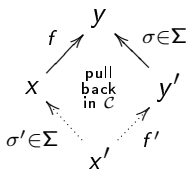


Yoneda system of a small category \mathcal{C} (2004-2007)

axiomatizing the preservation of the future and the past (2)

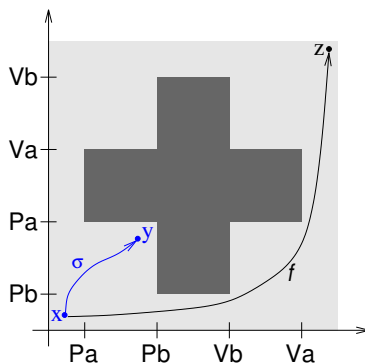
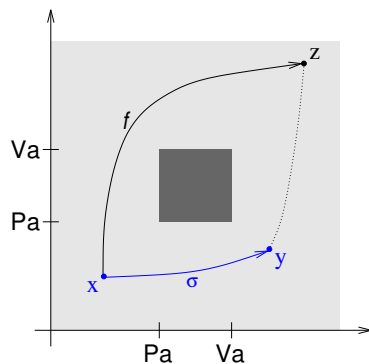
A collection Σ of morphisms of \mathcal{C} such that:

1. Σ is stable under composition,
2. Σ contains all the isomorphisms of \mathcal{C} ,
3. all the elements of Σ are *Yoneda* morphisms and
4. Σ is stable under **change** and **cochange** of base.



Examples

of morphisms which do not belong to a *Yoneda* system

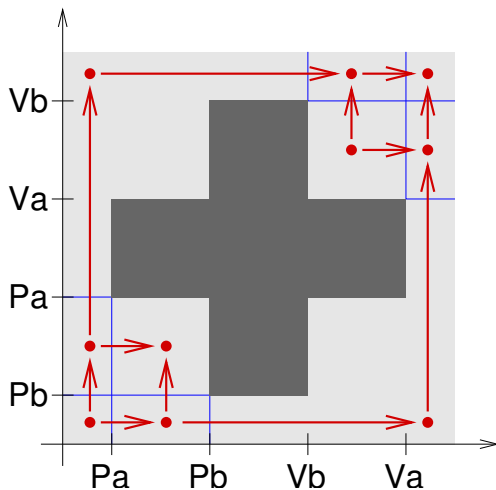


Fundamental properties

- ▶ Given Σ a Yoneda system, the isomorphism classes of objects (i.e. points of underlying X) is *convex*
 - ▶ Due to the *essential* property of **pureness**: for all $\sigma \in \Sigma$, $\sigma = f \circ g$ implies $f \in \Sigma$ and $g \in \Sigma$
- ▶ When C is **loop-free** i.e. **all isomorphisms are identities**, the collection of Yoneda systems forms a **locale**, i.e. is a complete lattice (infinite join and meet), which is distributive (and distributivity of meet over infinite join)

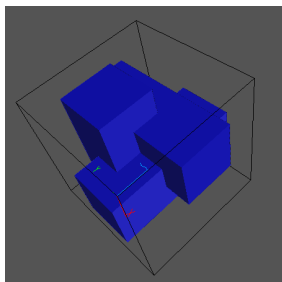
Define now the **component category** to be $\overrightarrow{\pi_0}(X)$ equal to the category of fractions of $\pi_1(X)$ by the maximal Yoneda System (equivalently as we shall see, as the quotient category of the same two categories).

The category of components of the Swiss flag

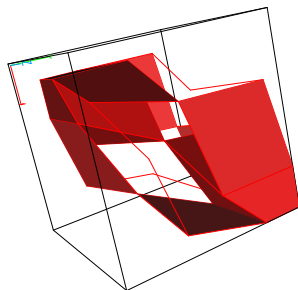


(the two red squares are commutative!)

The components category of the 3 philosophers

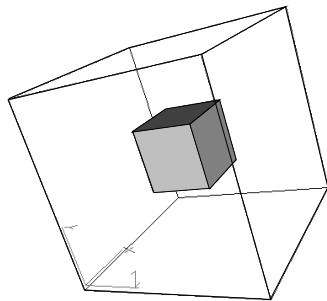


the pospace

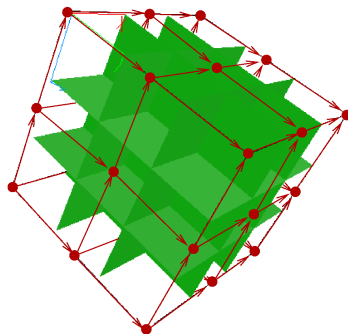


its category of components

The components category of a 2-semaphore



the pospace

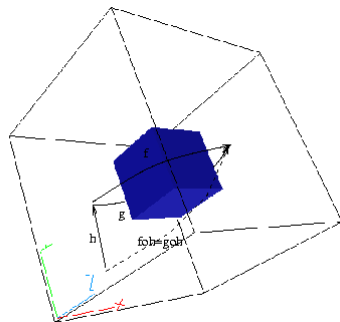


its category of components

(details of the calculation omitted...but...)

2 semaphore

Notice: a certain amount of the classical π_2 is apparent; and $\overrightarrow{\pi_1}$ has no “cancellation” property in general



Fundamental properties (2)

Lifting properties of the component category

Furthermore, let $C_1, C_2 \subset Ob(\mathcal{C})$ denote two components such that the set of morphisms (in \mathcal{C}/Σ) is *finite*. Then, for every $x_1 \in C_1$ there exists $x_2 \in C_2$ such that the **quotient map**

$$\begin{array}{ccc} \mathcal{C}(x_1, x_2) & \rightarrow & \mathcal{C}/\Sigma(C_1, C_2) \\ f & \mapsto & [f] \end{array}$$

is **bijjective**.

Fundamental results (3)

fractions vs quotients

Let \mathcal{C} be a small loop-free category and Σ a *Yoneda* system of \mathcal{C} :

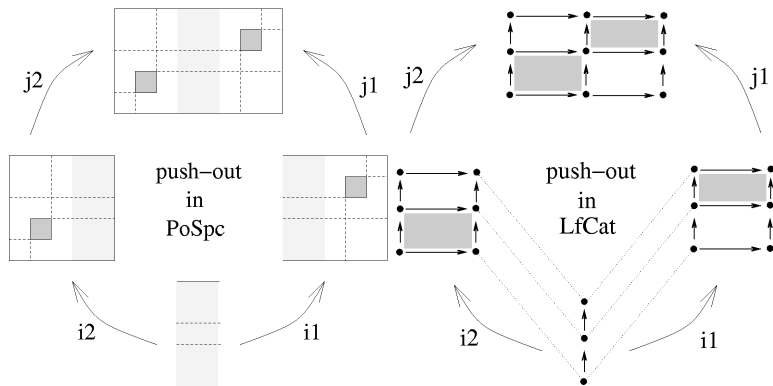
1. the small category \mathcal{C}/Σ is loop-free,
2. the small categories $\mathcal{C}[\Sigma^{-1}]$ and \mathcal{C}/Σ are equivalent and
3. the category $\mathcal{C}[\Sigma^{-1}]$ is fibered over \mathcal{C}/Σ .
4. **Seifert/van Kampen on component categories**
 - ▶ Relies heavily on the localic structure of Yoneda systems
 - ▶ Allows for inductive computations of $\mathcal{C}[\Sigma^{-1}]$ for PV programs
5. if \vec{K} is a compact pospace, then any component of $\vec{\pi}_1(\vec{K})$ has both a **greatest lower bound** and an **least upper bound** in $(|K|, \sqsubseteq)$.

see *Ph. D. Thesis* of E. Haucourt

see *Components of the Fundamental Category* - APCS 04, L. Fajstrup, E. Goubault, E. Haucourt, M. Raussen

see also *Components of the Fundamental Category II* - APCS 07, E. Goubault, E. Haucourt

Example of computation using van Kampen



→ inductive formulas in some cases, “algebra” of cubes etc. Used for static analysis (reduction of the state space)

Some figures - full component computations

(with the current naive implementation)

n	sec	Mb	# o	# m	# r	# p	#s	# t
3	0.38	≤ 10	27	48	18	6	576	1475
4	0.43	≤ 15	85	200	132	24	3966	13450
5	0.69	19	263	770	730	120	27265	113938
6	3.49	23	807	2832	3516	720	184876	914019
7	96	42	2467	10094	15484	5040	?	?
8	1656	100	7533	35216	64312	40320	?	?
9	13739	319	22995	120924	256158	362880	2996970	22698700

→ but non economical, need of future/past components only!

⇒ How can we generalize it to **future components**?

A closer view on components - “future” components (2008-)

Let \mathcal{C} be a small category, $\Sigma^+ \subseteq Mo(\mathcal{C})$ is a *Yoneda-f-system* if and only if Σ^+ is stable under composition and satisfies

(Af1) for all objects z of \mathcal{C} such that $\mathcal{C}[y, z] \neq \emptyset$, the map:

$$\mathcal{C}[y, z] \xrightarrow{-\circ\sigma} \mathcal{C}[x, z] \text{ is a bijection,$$

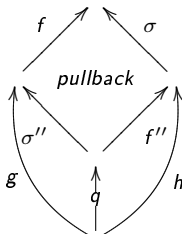
(Af2) Σ^+ contains $Iso(\mathcal{C})$

(Af3) Σ^+ is stable under pushouts (with any morphism in \mathcal{C})

And...

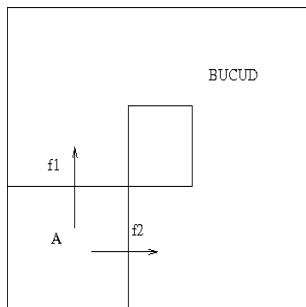
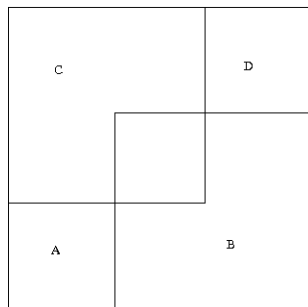
Future components

(Af4) Σ^+ is stable under filtered pullbacks (with any morphism in \mathcal{C}) i.e.: for all $f \in \mathcal{C}$, for all $\sigma \in \Sigma^+$ with same codomains such that there exists some $g \in \mathcal{C}$ and $h \in \mathcal{C}$ such that $f \circ g = \sigma \circ h$, then there exists $\sigma'' \in \Sigma^+$ giving us a pullback diagram:



Let $\mathcal{Y}_f(\mathcal{C})$ denote the collection of all Yoneda f -systems in \mathcal{C} .

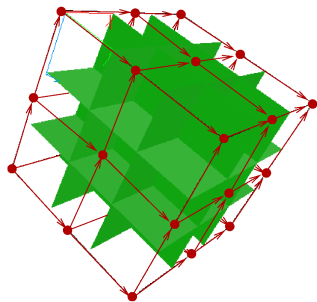
Example: $X = [0, 1]^2 \setminus]\frac{1}{3}, \frac{2}{3}[^2$



Notice that each of the component is a the trace on X of a **maxplus polyhedron**

Example: $X = [0, 1]^3 \setminus \frac{1}{3}, \frac{2}{3}[^3$

Start with the 26 below: (x, y, z) with $x, y, z \in \{-, 0, +\} \setminus \{(0, 0, 0)\}$ and the unique morphisms $f_{(xa, yb, zc)}$ in $\mathcal{C}((x, y, z), (a, b, c))$ with $x, y, z \in \{-, 0, +\} \setminus \{(0, 0, 0), (+, +, +)\}$ and (a, b, c) being $(s(x), y, z)$ or $(x, s(y), z)$ or $(x, y, s(z))$ (whenever defined), where $s(-) = 0$, $s(0) = +$ and $s(+)$ is undefined.



We know that Yoneda systems are particular future Yoneda systems, so we just have to check whether the $f_{(xa, xb, xc)}$ are in the maximal future Yoneda system or not:

Weakly-invertible morphisms in the future (WI)

- ▶ As $|\mathcal{C}((- , 0, -), (+, 0, +))| = 2$ whereas $|\mathcal{C}((- , - , -), (+, 0, +))| = 1$, morphisms in $\mathcal{C}((- , - , -), (-, 0, -))$ cannot be WI: $f_{(-, -, -), (-, 0, -)}$ = the class of such morphisms.
- ▶ By symmetry, morphisms in $\mathcal{C}((- , - , -), (0, - , -))$ and in $\mathcal{C}((- , - , -), (-, - , 0))$ cannot be WI: $f_{(-, -, -), (0, -, -)}$ and $f_{(-, -, -), (-, -, 0)}$ resp. these two morphisms.

Filtered pullback argument

- ▶ Consider $f_{(-0,00,--)}$ and $f_{(00,-0,--)}$:
 $f_{(-0,00,--)} \circ f_{(--,-0,--)} = f_{(00,-0,--)} \circ f_{(-0,--,--)}.$
- ▶ So if the morphism $f_{(-0,00,--)}$ were in Σ^+ then by (Af4), the diagram:

$$\begin{array}{ccc}
 & f_{(-0,00,--)} & \\
 & \dashrightarrow & \\
 f_{(--,-0,--)} & \uparrow & f_{(00,-0,--)} \\
 & \dashrightarrow & \\
 & f_{(-0,--,--)} &
 \end{array}$$

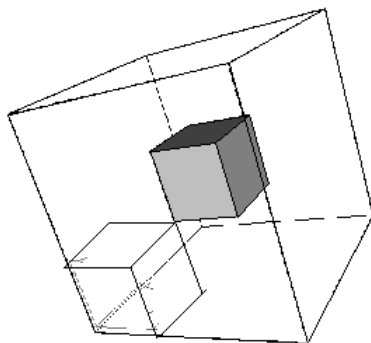
would have to be a pullback diagram with $f_{(-0,--,--)}$ in Σ^+ .

- ▶ But $f_{(-0,--,--)}$ is not WI! so impossible!
- ▶ Hence $f_{(-0,00,--)}$ is not in Σ^+ , similarly for $f_{(00,-0,--)}$

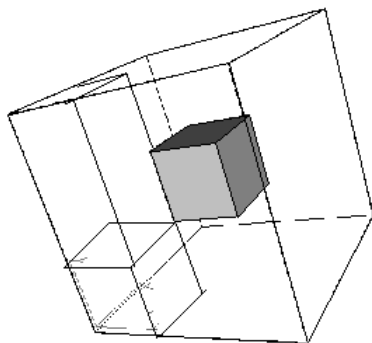
Iteration of the argument...

- ▶ By symmetry, this argument shows that $f_{(-0,--,00)}$, $f_{(00,00,-0)}$, $f_{(--,-0,00)}$, and $f_{(--,00,-0)}$ are not in Σ^+ .
- ▶ And then... $f_{(-0,yy,zz)}$, $f_{(xx,-0,zz)}$ and $f_{(xx,yy,-0)}$ with $x, y, z \in \{-, 0, +\}$ cannot be in Σ^+ (by “propagation” around one axis of the filtered pullback argument above, of each of the three morphisms that are not future Yoneda invertible)
- ▶ No other constraints can be found due to axioms (Af1) to (Af4).

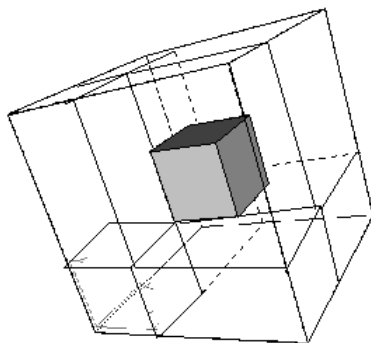
Future components, geometrically



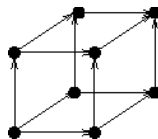
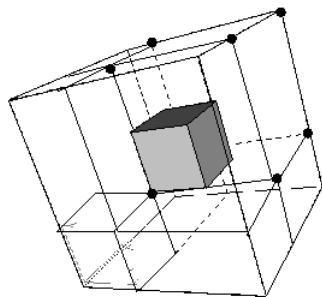
Future components, geometrically



Future components, geometrically



i.e....



\implies Take **full subcat** of $\vec{\pi}_1(\vec{X})$ with “**max of components**” objects!

We will give meaning to this later

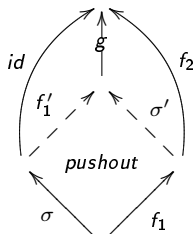
Certainly not true with full components...

Fundamental results (1): Yoneda- f -systems are pure

Suppose $Iso(\mathcal{C})$ is pure, then any Yoneda- f -system is pure.

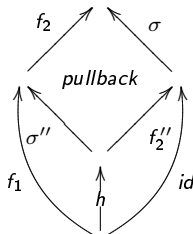
Proof.

- ▶ Take $\Sigma^+ \in \mathcal{Y}_f(\mathcal{C})$ and take $\sigma \in \Sigma^+$ and $f_1, f_2 \in Mo(\mathcal{C})$ such that $\sigma = f_2 \circ f_1$.
- ▶ By (Af3), we have a $\sigma' \in \Sigma^+$ and f_1' and a unique $g \in Mo(\mathcal{C})$ s.t.:



Proof...

- ▶ By pureness of $Iso(\mathcal{C})$ in \mathcal{C} , f_1' and g are isomorphisms
- ▶ hence by (Af2), belong to Σ^+ .
- ▶ So by (Af1), $f_2 = g \circ \sigma' \in \Sigma^+$.
- ▶ Now, f_2 and $\sigma \in \Sigma^+$ admit a pullback by (Af4) because $f_2 \circ f_1 = \sigma \circ id$:



where σ'' is in Σ^+ .

End of proof

- ▶ Also, $id = f_2'' \circ h$ and as $Iso(\mathcal{C})$ is pure, h is in $Iso(\mathcal{C})$
- ▶ By (Af2), h is in Σ^+ .
- ▶ We also have in the commutative diagram of last slide that $f_1 = \sigma'' \circ h$, composite of two arrows in Σ^+ , hence in Σ^+
- ▶ Thus Σ^+ is pure in \mathcal{C} .

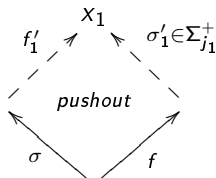


The locale of Yoneda-f-systems (meet)

If $(\Sigma_j^+)_{j \in J}$ is a non empty family of *Yoneda-f-systems* of a small category \mathcal{C} then $\bigcap_{j \in J} \Sigma_j^+$ is a *Yoneda-f-system* of \mathcal{C} .

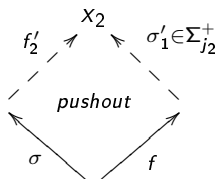
Proof.

- ▶ (Af1) and (Af2) are trivial
- ▶ Suppose $\sigma \in \bigcap_{j \in J} \Sigma_j^+$ and $f \in Mo(\mathcal{C})$ with $src(f) = src(\sigma)$.
Take $j_1, j_2 \in J$, since $\sigma \in \Sigma_{j_1}^+$ we have a pushout square



The locale of Yoneda-f-systems (meet)

- ▶ and also

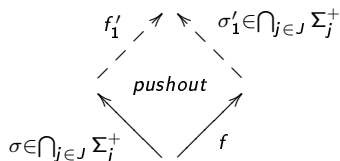


because $\sigma \in \Sigma_{j_2}^+$

- ▶ By uniqueness of the pushout, we have an iso τ from x_2 to x_1 s.t. $\sigma'_1 = \tau \circ \sigma'_2$. (Af2) implies $\tau \in \Sigma_{j_2}^+$ and (Af1) implies $\sigma'_1 = \tau \circ \sigma'_2 \in \Sigma_{j_2}^+$

The locale of Yoneda-f-systems (meet)

- By the same argument, for all $j \in J, \sigma'_1 \in \Sigma_j^+$ i.e. $\sigma'_1 \in \bigcap_{j \in J} \Sigma_j^+$ and we have



- The same proof holds for pullback squares ■

The locale of Yoneda-f-systems (join)

If $(\Sigma_j^+)_{j \in J}$ is a non empty family of *Yoneda*-f-systems of a small category \mathcal{C} then $\biguplus_{j \in J} \Sigma_j^+$ is a *Yoneda*-system of \mathcal{C} , where $\biguplus_{j \in J} \Sigma_j^+$ is the least sub-category of \mathcal{C} including all the Σ_j^+ 's.

Proof.

- ▶ $\biguplus_{j \in J} \Sigma_j^+ = \{\sigma_n \circ \dots \circ \sigma_1 \mid n \in \mathbb{N}^*, \{j_1, \dots, j_n\} \subseteq J \text{ and for all } k \in \{1, \dots, n\}, \sigma_k \in \Sigma_{j_k}^+\}, \implies (\text{Af1})$
- ▶ (Af2) also trivial

□

The locale of Yoneda-f-systems (join)

- ▶ Take $\sigma_n \circ \dots \circ \sigma_1 \in \bigoplus_{j \in J} \Sigma_j^+$ with $n \in \mathbb{N}^*$, $\{j_1, \dots, j_n\} \subseteq J$, for all $k \in \{1, \dots, n\}$, $\sigma_k \in \Sigma_{j_k}^+$ and $f \in Mo(\mathcal{C})$ with $src(\sigma_1) = src(f)$:
- ▶

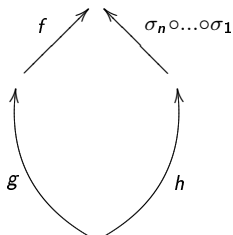
$$\begin{array}{c}
 \uparrow f \\
 \xrightarrow{\sigma_1 \in \Sigma_{j_1}^+} \dots \xrightarrow{\sigma_n \in \Sigma_{j_n}^+}
 \end{array}$$

By a finite induction using (Af3) for $\Sigma_{j_1}^+, \dots, \Sigma_{j_n}^+$:

$$\begin{array}{ccc}
 \sigma'_1 \in \Sigma_{j_1}^+ & & \sigma'_n \in \Sigma_{j_n}^+ \\
 \xrightarrow{\quad} & & \xrightarrow{\quad} \\
 \uparrow & \wedge & \uparrow \\
 f \uparrow & p.o. \mid f_1 & f_{n-1} \mid p.o. \mid f_n \\
 \xrightarrow{\sigma_1 \in \Sigma_{j_1}^+} & \dots & \xrightarrow{\sigma_n \in \Sigma_{j_n}^+} \\
 \uparrow & & \uparrow
 \end{array}$$

The locale of Yoneda-f-systems (join)

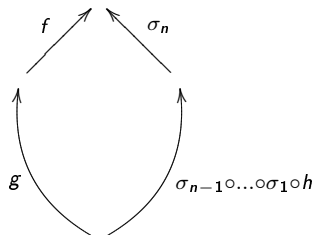
- Suppose



with $\sigma_i \in \Sigma_{u(i)}^+$ for all $i = 1, \dots, n$ and u is a function from $\{1, \dots, n\}$ to I

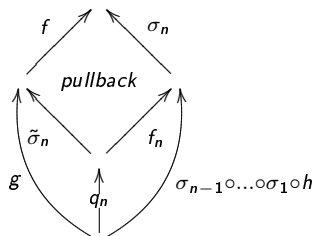
The locale of Yoneda-f-systems (join)

► So



The locale of Yoneda-f-systems (join)

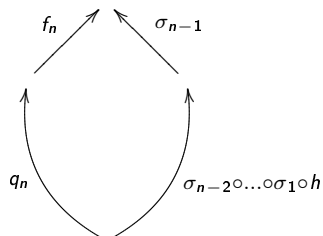
- (Af4) for $\Sigma_{u(n)}^+$ implies:



with $q_n, f_n \in \mathcal{C}$ and $\tilde{\sigma}_n \in \Sigma_{u(n)}^+$.

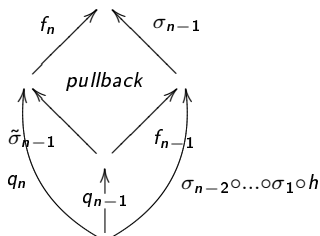
The locale of Yoneda-f-systems (join)

- Now we carry on with:



The locale of Yoneda-f-systems (join)

- hence by the same argument: diagram:



with $q_{n-1}, f_{n-1} \in \mathcal{C}$ and $\tilde{\sigma}_{n-1} \in \Sigma_{u(n-1)}^+$.

The locale of Yoneda-f-systems (join)

- By induction: of n pullback squares:

$$\begin{array}{ccc}
 \tilde{\sigma}_1 \in \Sigma_{u(1)}^+ & & \tilde{\sigma}_n \in \Sigma_{u(n)}^+ \\
 \begin{array}{ccc}
 \dashrightarrow & \dashrightarrow & \\
 \downarrow f_1 & \text{p.b.} & \downarrow f_2 \\
 \dashrightarrow & & \dashrightarrow \\
 \sigma_1 \in \Sigma_{u(1)}^+ & & \sigma_n \in \Sigma_{u(n)}^+
 \end{array} & \cdots & \begin{array}{ccc}
 \dashrightarrow & \dashrightarrow & \\
 \downarrow f_n & \text{p.b.} & \downarrow f \\
 \dashrightarrow & & \dashrightarrow \\
 \sigma_n \in \Sigma_{u(n)}^+ & & \sigma_n \in \Sigma_{u(n)}^+
 \end{array}
 \end{array}$$

with $\tilde{\sigma}_n \circ \tilde{\sigma}_1 \in \Sigma^+$



Conclusion... (skipping some lemmas)

- ▶ Let \mathcal{C} be a small category such that $Iso(\mathcal{C})$ is pure in \mathcal{C} . Then,
 - ▶ the family of future Yoneda-systems of \mathcal{C} is not empty and,
 - ▶ together with \subseteq it forms a locale
 - ▶ whose l.u.b. operator is \biguplus and g.l.b operator is \bigcap .
 - ▶ Moreover, the least element of this locale (“bottom”) is $Iso(\mathcal{C})$
- ▶ There exists a maximal future Yoneda system in any category \mathcal{C} .

\implies Notion of future component category $\overrightarrow{\pi}_0(X)$

It all works as with “full” components...

- ▶ We have a **van Kampen** theorem on future components (consequence of the locale structure)
 - ▶ For PV programs, allows for proving components are maxplus polyhedra of some sort; inductive calculation and “algebra” of maxplus polyhedra
- ▶ We have an equivalence between category of fractions and quotient category
- ▶ We have a **lifting property** (slightly different - only one side!): For every $x_1 \in C_1$ there exists $x_2 \in C_2$ such that the quotient map

$$\mathcal{C}(x_1, x_2) \rightarrow \mathcal{C}/\Sigma^+(C_1, C_2), f \mapsto [f]$$

is bijjective.

An extra condition on components (?)

(Conjecture(?)): automatically true in the PV case

- ▶ Let \mathcal{D} be the category whose objects are $X, X_0, X_1, \dots, X_n, \dots$, and whose only morphisms are of the form $X \rightarrow X_i$ ($i \geq 0$).
- ▶ Let F be a functor from \mathcal{D} to a category \mathcal{C} .
- ▶ We call infinite pushout the colimit of $F(\mathcal{D})$ in \mathcal{C} , when it exists.

Ask for future (resp. past) components to have **infinite pushouts** (resp. pullbacks).

(They already had finite pushouts)

Extension of the lifting property

With this extra property, we have both for past and future components:

- ▶ the lifting property holds
- ▶ even if the set of morphisms (in \mathcal{C}/Σ) between two objects is not finite.

Orthogonal subcategories

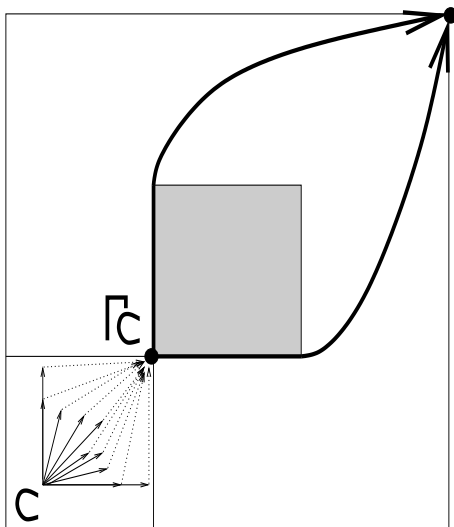
See e.g. Borceux

Let \mathcal{C} be a category and Σ a class of morphisms of \mathcal{C} .

- ▶ By the orthogonal subcategory of \mathcal{C} determined by Σ , we mean the full subcategory \mathcal{C}_Σ of \mathcal{C} ,
- ▶ whose objects are those $X \in \mathcal{C}$ such that $s \perp X$ for every $s \in \Sigma$, i.e.,
- ▶ such that for every $s : A \rightarrow B \in \Sigma$, for every morphism $f : A \rightarrow X$, there exists a unique morphism $b : B \rightarrow X$ such that $b \circ s = f$.

$$\begin{array}{ccc}
 A & \xrightarrow{s \in \Sigma} & B \\
 \downarrow \forall f \in \mathcal{C} & \swarrow \exists! b & \\
 X & &
 \end{array}$$

The orthogonal subcategory of Σ_+ is reflective



Theorem

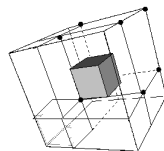
Let Σ be the inessential morphisms in the future, in the category $\mathcal{C} = \vec{\pi}_1(\vec{X})$ for some local po-space X .

Suppose that Σ has infinite pushouts then

- ▶ \mathcal{C}_Σ is equivalent to $\mathcal{C}[\Sigma^{-1}]$

- ▶ \mathcal{C}_Σ is reflective in $\vec{\pi}_1(\vec{X})$

(note that $\vec{\pi}_1(\vec{X})$ is in general not complete, and that not all objects are representable!)



This gives (indirectly) a reason why we had:

Sketch of proof (2nd part)

By definition of the orthogonal subcategory:

- ▶ we have an obvious inclusion functor I from \mathcal{C}_Σ to \mathcal{C} .
- ▶ To prove that we have a reflective subcategory, we need to construct the left adjoint to $\Gamma \dashv I$.

Sketch of proof

- ▶ Let $C \in \mathcal{C}$. For every pair (s, f) where $s : S \rightarrow T \in \Sigma$ and $f : S \rightarrow C \in \mathcal{C}$,
- ▶ by the properties of Σ , we know we have a pushout diagram (we call it a (s, f) pushout square):

$$\begin{array}{ccc}
 S & \xrightarrow{s \in \Sigma} & T \\
 f \downarrow & & \downarrow g_{sf} \\
 C & \xrightarrow[t_{sf} \in \Sigma]{} & P_{sf}
 \end{array}$$

where $t_{sf} \in \Sigma$.

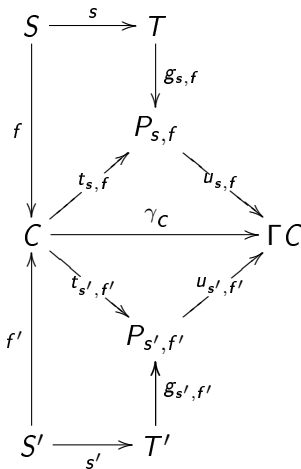
Sketch of proof

Consider now the diagram composed of all t_{sf} for all pairs (s, f) which is small since Σ is a set, thus its colimit $(\Gamma C, (u_{sf})_{sf})$ exists in $\underline{\Sigma}$ and provide $\Gamma C \in \Sigma$.

We have defined the object part of Γ . Now we construct a family of morphisms of \mathcal{C} , denoted $(\gamma_C)_{C \in \text{Ob } \mathcal{C}}$, that will be the unit of the adjunction.

Sketch of a proof

Let us determine $\gamma_C : C \rightarrow \Gamma C$. Given two (s, f) -pushout squares, by definition of a colimit, we have $u_{s,f} \circ t_{s,f} = u_{s',f'} \circ t_{s',f'}$, hence we can set $\gamma_C := u_{s,f} \circ t_{s,f}$ since it does not depend on the (s, f) -pushout square we have chosen. $\gamma_C \in \Sigma$ for it is given by the composite of two morphisms of Σ .



Sketch of proof

We determine the morphism part of Γ , this construction will implicitly prove that γ is an adjunction from Id_C to $I \circ \Gamma$. Let $h : C^1 \rightarrow C^2$. For each pushout square

$$\begin{array}{ccc} S & \xrightarrow{s} & T \\ f \downarrow & & \downarrow g_{s,f}^1 \\ C^1 & \xrightarrow{t_{s,f}^1} & P_{s,f}^1 \end{array}$$

Sketch of proof

We have the pushout diagram (since $t_{s,f}^1 \in \Sigma$):

$$\begin{array}{ccc} C^1 & \xrightarrow{t_{s,f}^1} & P_{s,f}^1 \\ \downarrow h & & \downarrow g_{t_{s,f}^1}^2 \\ C^2 & \xrightarrow{t_{s,f,h}^2} & P_{s,f,h}^2 \end{array}$$

Sketch of proof

Then, using the factorisation of the unit (on C^1 and C^2) we have:

$$\begin{array}{ccccc}
 C^1 & \xrightarrow{t_{s,f}^1} & P_{s,f}^1 & \xrightarrow{u_{s,f}^1} & \Gamma(C^1) \\
 \downarrow h & & \downarrow g_{t_{s,f}^1, h}^2 & & \\
 C^2 & \xrightarrow{t_{t_{s,f}^1, h}^2} & P_{t_{s,f}^1, h}^2 & \xrightarrow{u_{t_{s,f}^1, h}^2} & \Gamma(C^2)
 \end{array}$$

Sketch of proof

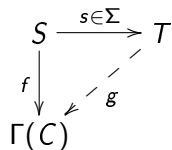
Since $u_{s,f}^1$ is in Σ , it is invertible in the future, giving the map $\Gamma(h)$ making the following diagram commute:

$$\begin{array}{ccccc}
 C^1 & \xrightarrow{t_{s,f}^1} & P_{s,f}^1 & \xrightarrow{u_{s,f}^1} & \Gamma(C^1) \\
 \downarrow h & & \downarrow g_{t_{s,f}^1}^2 & & \downarrow \Gamma(h) \\
 C^2 & \xrightarrow{t_{s,f,h}^2} & P_{t_{s,f,h}^1}^2 & \xrightarrow{u_{t_{s,f,h}^1}^2} & \Gamma(C^2)
 \end{array}$$

Sketch of proof

Now we prove that $\forall C \in \mathcal{C}, \Gamma(C) \in \mathcal{C}_\Sigma$.

Given (s, f) with $s : S \rightarrow T \in \Sigma$, $src(f) = S$ and $tgt(f) = \Gamma(C)$, we have a g making the right side diagram commutative. Precisely, $g := u_{sf} \circ g_{sf}$, indeed, $u_{sf} \circ g_{sf} \circ s = u_{sf} \circ t_{sf} \circ f = \gamma_C \circ f$.



Sketch of proof

- ▶ The uniqueness of g is due to the bijectivity of $\gamma \in \mathcal{C}[S, T] \longrightarrow \gamma \circ s \in \mathcal{C}[S, \Gamma(C)]$, because s is inessential in the future.
- ▶ Thus, $s \perp \Gamma(C)$ and $\Gamma(C)$ is in the orthogonal subcategory determined by Σ .

Sketch of proof

Conversely, suppose that $X \in \mathcal{C}_\Sigma$. Given (s, f) with $s : S \rightarrow T \in \Sigma$, $\text{src}(f) = S$ and $\text{tgt}(f) = X$, we have a unique g making the right side diagram commutative, which is in fact necessarily the pushout square given by the definition of Σ (since s is epi due to (Af2)).

$$\begin{array}{ccc}
 S & \xrightarrow{s \in \Sigma} & T \\
 f \downarrow & & \downarrow g \\
 X & \xrightarrow{id_X} & X
 \end{array}$$

Sketch of proof

With the notation introduced at the beginning of the proof,

- ▶ $g_{s,f} = g$ and $t_{s,f} = id_X$.
- ▶ But then the colimit of such pushout squares, defining $(\Gamma(X), u_{s,f})$ is the colimit of the family $(id_X)_{\{(s,f) \text{ with } s:S \rightarrow T \in \Sigma, \text{ src}(f)=S \text{ and } \text{tgt}(f)=X\}}$.
- ▶ Hence $u_{s,f} \cong id_X$ for all such pairs (s, f) and $\Gamma(X) \cong X$ in \mathcal{C} .

Sketch of proof

The last part of the proof consists in seeing that

$$\alpha \in \mathcal{C}_\Sigma[\Gamma(C), D] \longmapsto I(\alpha) \circ \gamma_C \in \mathcal{C}[C, I(D)]$$

is a bijection, where $\gamma_C : C \rightarrow \Gamma C$ is the canonical morphism given by the colimit.

Sketch of proof

Given $D \in \mathcal{C}_\Sigma$ and $m : C \rightarrow D$, we must find a unique $n : \Gamma C \rightarrow D$ such that $n \circ \gamma_C = m$.

- ▶ As $D \in \mathcal{C}_\Sigma$, it is orthogonal to all morphisms $s \in \Sigma$.
- ▶ In particular, for all pairs (s, f) as above, there exists a unique b_{sf} such that $b_{sf} \circ s = m \circ f$.
- ▶ By the pushout property defining P_{sf} we deduce that there is a unique morphism $a_{sf} : P_{sf} \rightarrow D$ such that $a_{sf} \circ t_{sf} = m$ and $a_{sf} \circ g_{sf} = b_{sf}$.

Sketch of proof

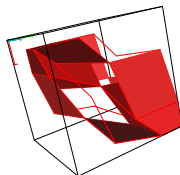
- ▶ This is done for all pairs (s, f) .
- ▶ Hence by the colimit property defining ΓC , we find a unique morphism $n : \Gamma C \rightarrow D$ such that $n \circ u_{sf} = \alpha_{sf}$.
- ▶ Hence

$$\begin{aligned}n \circ \gamma_C &= n \circ u_{sf} \circ t_{sf} \\ &= \alpha_{sf} \circ t_{sf} \\ &= m\end{aligned}$$

which ends the proof.

Mathematical improvements

- ▶ Computation of morphisms in the component category:
 - ▶ Help from **homology**? (many tries!) Promising one: non-abelian cohomology (using coefficients in semi-rings)
 - ▶ (on the longer run) simplification of the retract by considering also the $\vec{\pi}_n$ [“homotopical resolution”, or “**higher-order syzygies**”], example:

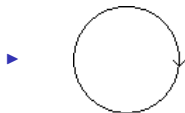


1st order: $[a,b],$
 $[b,c], [a,c]$
2nd order: $[a,b,c]$

- ▶ Use of **trace spaces**? (M. Raussen)

Open mathematical issues

- ▶ Components in the presence of **loops/non-determinism**:



the category of components is exactly π_1
(continuum of points)!

difficult to find a direct notion of Yoneda
system in that case

we would like something like
components $\sim (\mathbb{N}, +)$

- ▶ Should be coherent with the view that one would have when looking at the **universal discovering** (see recent work by L. Fajstrup, also S. Krishnan)

Open mathematical issues

- ▶ **dihomotopy equivalence?** Algebraic homotopical structure (Quillen, Baues...)?
⇒ leads in work by M. Grandis, P. Gaucher, K. Worytkiewicz etc.
- ▶ **Classification** of models modulo dihomotopy equivalence? (at least in dimension 2?)
⇒ right “**primitive**” **synchronisation primitives?** (foundational work, concurrent Turing machine?)
- ▶ Enrichment of the model, e.g. timing issues:
⇒ geodesics, CAT0 conditions (see R. Ghrist/V. Paterson)

Open mathematical issues

- ▶ Right directed notions in **cubical sets**?
⇒ would be much more natural in the context of semantics and static analysis - leads in work (classical case) of R. Jardine
 - ▶ What is the correct category involved?
 - ▶ Can we define directly an algebraic homotopical structure (for the directed equivalence)? Quillen equivalence with the (di-)topological case?
- ▶ Relation with **causality** issues in theoretical physics (general relativity)?
⇒ see Penrose models, P. Panangaden/K. Martin work
- ▶ Relationship with **topological complexity** (Mike Farber)? with classical work in **combinatorial algebraic topology** (arrangement spaces, see Dimitri Kozlov)? ...

Thanks for your attention!

Related Work

- ▶ “Components of the Fundamental Category II”, E. Goubault, E. Haucourt, in Applied Categorical Structures 07
- ▶ “Components of the Fundamental Category”, L. Fajstrup, E. Goubault, E. Haucourt, M. Raussen, in Applied Categorical Structures 04.
- ▶ “Algebraic Topology and Concurrency ”, L. Fajstrup, E. Goubault, M. Raussen, (MFPS'98) Theoretical Computer Science 05.
- ▶ “Detecting Deadlocks in Concurrent Systems ”, L. Fajstrup, E. Goubault, M. Raussen (Concur'98).
- ▶ “Some Geometric Perspectives in Concurrency Theory ”, E. Goubault, in Homology, Homotopy and Applications 03.

Related Work

- ▶ M. Raussen and L. Fajstrup (many things, among which a complete treatment for dimension 2, characterization of obstructions to dihomotopies etc.)
- ▶ M. Grandis “quotient models” and higher fundamental categories (a lot there again)
- ▶ R. Ghrist and V. Peterson (only mutex, but geodesics!)
- ▶ (partial order reduction methods, Godefroid, Wolper...) Ample sets, persistent sets, stubborn sets, sleep sets etc.