

The time constant for Bernoulli percolation is Lipschitz continuous strictly above p_c

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Percolation

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- Graph $(\mathbb{Z}^d, \mathbb{E}^d)$, $d \geq 2$.
- $(B(e))_{e \in \mathbb{E}^d}$: i.i.d. family of Bernoulli random variable of parameter $p \in [0, 1]$.
- $B(e) = 1 \implies e$ is open.
- $B(e) = 0 \implies e$ is closed.

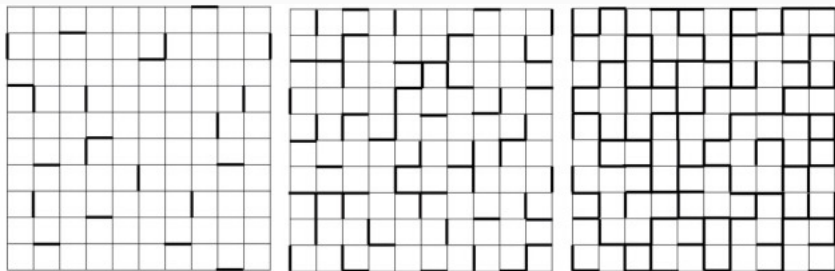


Figure 1: Simulation of percolation for parameters $p = 0.1$; 0.3 and 0.6

Percolation probability

- Random graph $\mathcal{G}_p = (\mathbb{Z}^d, \{e \in \mathbb{E}^d : B(e) = 1\})$.
- $\mathcal{C}_p(0)$: the connected component of 0 in \mathcal{G}_p .

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- $\theta(0) = 0$
- $\theta(1) = 1$
- $p \mapsto \theta(p)$ is nondecreasing

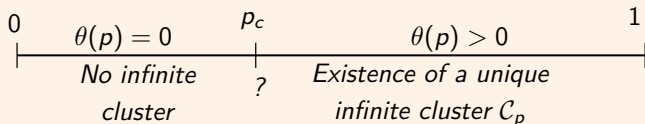
Phase transition

Definition (Critical parameter)

$$p_c = \sup \{ p : \theta(p) = 0 \}$$

Phase transition at $p_c \in]0, 1[$:

Theorem (Broadbent-Hammersley 57-59, ...)



First passage percolation

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What is the asymptotic value of $T_G(0, nx)$?

Theorem (Hammersley-Welsh 65, Kingman 73-75)

Under some conditions on G , we have

$$\forall x \in \mathbb{Z}^d \quad \lim_{n \rightarrow \infty} \frac{T_G(0, nx)}{n} = \mu_G(x) \text{ almost surely and in } L^1.$$

where $\mu_G(x)$ is a deterministic constant. This is the so-called time constant.

First passage percolation : Definition of the time constant

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Regularity of μ_G in G ?

Theorem (Cox 81, . . . , Garet-Marchand-Proccacia-Théret 17)

The map $G \mapsto \mu_G$ is continuous.

Time constant in the Bernoulli case

Graph distance

We are interested in the random metric induced by \mathcal{G}_p when $p > p_c$. We define for x and y in \mathbb{Z}^d

$$\mathcal{D}_p(x, y) = \inf \{ |\gamma| : \gamma \text{ path that joins } x \text{ and } y \text{ in } \mathcal{G}_p \}$$

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Set $G_p = p\delta_1 + (1-p)\delta_\infty$. We can couple $(t(e))_{e \in \mathbb{E}^d}$ with $(B(e))_{e \in \mathbb{E}^d}$ by setting $B(e) = \mathbb{1}_{t(e)=1}$ so that

$$\mathcal{D}_p = T_{G_p}.$$

First passage percolation : Definition of the time constant for the graph distance

Theorem (Cerf-Théret 14)

For $p > p_c$, for any $x \in \mathbb{Z}^d$, there exists $\mu_p(x) > 0$ such that

$$\lim_{n \rightarrow \infty} \frac{\mathcal{D}_p(\tilde{0}, \tilde{nx})}{n} = \mu_p(x) \text{ almost surely and in } L^1$$

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Theorem (D. 18)

Let $p_0 > p_c$, there exists a positive constant C (depending on p_0) such that

$$\forall p, q \in [p_0, 1] \quad \sup_{\|x\|=1} |\mu_p(x) - \mu_q(x)| \leq C|q - p| \log |q - p|.$$

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General idea of the proof

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$$\mathbb{P}(e \text{ is } p\text{-closed} \mid e \text{ is } q\text{-open}) = \mathbb{P}(U(e) \geq p \mid U(e) \leq q) = \frac{q-p}{q}$$

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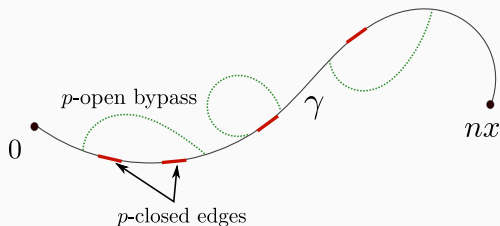


Figure 2: Build a p -open path upon a q -open path for $q > p > p_c$

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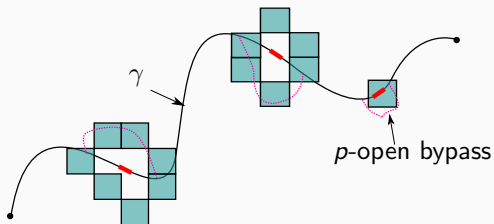
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$$\mathcal{D}_p(0, nx) \leq |\gamma'| \leq |\gamma| + |\gamma' \setminus \gamma| \leq \mathcal{D}_q(0, nx) + |\gamma' \setminus \gamma|$$

If we prove that $|\gamma' \setminus \gamma| \leq C_0|q - p|n$ then

$$\mu_p \leq \mu_q + C_0|q - p|.$$

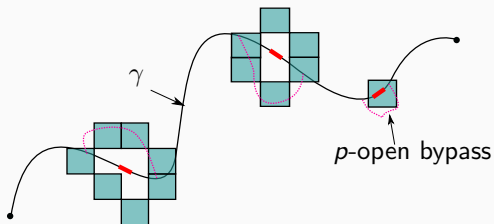
First approach: renormalization



- a good N -box
- a p -closed edge

Divide the lattice into boxes of mesoscopic size N . A good box is a box that has good connectivity property. Being a good box is something very likely for N large.

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Divide the lattice into boxes of mesoscopic size N . A good box is a box that has good connectivity property. Being a good box is something very likely for N large. Two cases :

1. Bad edge in good box
2. Bad edge in bad box

A different approach

Let $q > p > p_c$. γ is the q -geodesic between 0 and nx . We don't reveal which edges need to be bypassed. For each $e \in \gamma$, we define $c(e)$ the cost to bypass e such that:

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We have

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To build $c(e)$ we need a multiscale renormalisation. We need to consider an infinite number of scales of box at the same time. A good box at scale $k + 1$ is good if it does not contain too much bad boxes at scale k .

Thank you for your attention !